HIGHER-ORDER REWRITING SYSTEMS, CATEGORIAL ALGEBRAS, AND CURRY-HOWARD ISOMORPHISMS

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REFERENCES

J. Climent Vidal, E. Cosme Llópez. **From higher-order rewriting systems to higher-order categorial algebras and higher-order Curry-Howard isomorphisms**. ArXiv, abs/2402.12051, 2024.

REWRITING SYSTEM

A **rewriting system** is an ordered tuple $\boldsymbol{\mathcal{A}} = (\Sigma, X, \mathcal{A})$ where

- Σ is a signature;
- X is a set of variables;
- $\mathcal{A}$ is a subset of $\mathrm{T}_\Sigma(X)^2$.

The elements of $\mathcal{A}$ are called **rewriting rules**.

PATHS

A **path in $\mathcal{A}$** of length $m \in \mathbb{N}$ is

$$\mathfrak{P} = ((P_i)_{i \in m+1}, (\mathfrak{p}_i)_{i \in m}, (T_i)_{i \in m})$$

where, for every $i \in m$, if $\mathfrak{p}_i = (M_i, N_i)$, then

(1) $T_i(M_i) = P_i$; (2) $T_i(N_i) = P_{i+1}$.

$$\mathfrak{P}\colon\ P_0 \xrightarrow{(\mathfrak{p}_0, T_0)} P_1 \xrightarrow{(\mathfrak{p}_1, T_1)} \cdots \xrightarrow{(\mathfrak{p}_{m-2}, T_{m-2})} P_{m-1} \xrightarrow{(\mathfrak{p}_{m-1}, T_{m-1})} P_m$$

PATHS

Example

$$
\begin{aligned}
\mathfrak{P}:\ \oplus(x, \oplus(x, y)) &\xrightarrow{((y,z),\ \oplus(x,\oplus(x,\_)))} \oplus(x, \oplus(x, z)) \\
&\xrightarrow{((\oplus(x,z),z),\ \oplus(x,\_))} \oplus(x, z) \\
&\xrightarrow{((\oplus(x,z),\odot(\boxdot(z,x),z,\boxdot(x,x))),\ \_)} \odot(\boxdot(z, x), z, \boxdot(x, x)) \\
&\xrightarrow{((z,x),\ \odot(\boxdot(z,x),\_,\boxdot(x,x)))} \odot(\boxdot(z, x), x, \boxdot(x, x)) \\
&\xrightarrow{((\boxdot(z,x),y),\ \odot(\_,x,\boxdot(x,x)))} \odot(y, x, \boxdot(x, x)) \\
&\xrightarrow{((\boxdot(x,x),z),\ \odot(y,x,\_))} \odot(y, x, z) \\
&\xrightarrow{((\odot(y,x,z),\top),\ \_)} \top
\end{aligned}
$$

MAIN QUESTION

When can two rewriting systems be considered **equivalent**?

COMPOSITION

Paths can be composed.

If $\mathfrak{P}\colon P \longrightarrow Q$ and $\mathfrak{Q}\colon Q \longrightarrow R$, then $\mathfrak{Q} \circ \mathfrak{P}\colon P \longrightarrow R$.

Composition is a partial binary operation.

$$\mathrm{Pth}_{\mathcal{A}} \; \begin{array}{c} \xrightarrow{\ \mathrm{sc}\ } \\ \xleftarrow{\ \mathrm{ip}\ } \\ \xrightarrow[\ \mathrm{tg}\ ]{} \end{array} \; \mathrm{T}_{\Sigma}(X)$$

We denote by $\mathbf{Pth}_{\mathcal{A}}$ to the category whose objects are terms and whose morphisms are paths.

DECOMPOSITION

Paths can be decomposed.

If $\mathfrak{p} = (M, N)$ is a rewriting rule in $\mathcal{A}$, its associated **echelon** is the path of length 1

$$\mathrm{Ech}(\mathfrak{p}) : M \xrightarrow{(\mathfrak{p},\_)} N$$

We will say that a path has echelons if any of its subpaths of length 1 is an echelon.

DECOMPOSITION

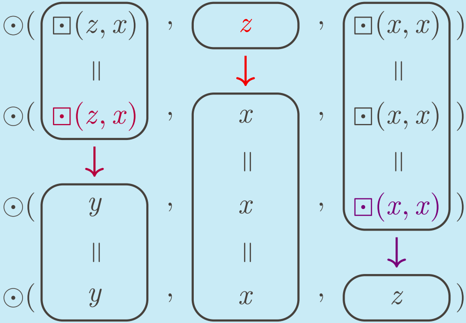
Example

$$
\begin{array}{rcl}
\mathfrak{P}: \oplus(x, \oplus(x, y)) & \to & \oplus(x, \oplus(x, z)) \\
 & \to & \oplus(x, z) \\
\text{echelon} & \to & \odot(\boxdot(z, x), z, \boxdot(x, x)) \\
 & \to & \odot(\boxdot(z, x), x, \boxdot(x, x)) \\
 & \to & \odot(y, x, \boxdot(x, x)) \\
 & \to & \odot(y, x, z) \\
\text{echelon} & \to & \top
\end{array}
$$

Proposition. Paths without echelons are paths between complex and homogeneous terms.

DECOMPOSITION

Example



Proposition. In a path without echelons, we can extract as many subpaths as the arity of the operation.

DECOMPOSITION

Let $\prec$ be the binary relation on $\mathrm{Pth}_{\mathcal{A}}$ defined by $\mathfrak{Q} \prec \mathfrak{P}$ if

i. $\mathfrak{P}$ has length strictly greater than 1, has its first echelon in position i and $\mathfrak{Q}$ is the prefix subpath strictly preceding the echelon or the suffix subpath containing the echelon; or

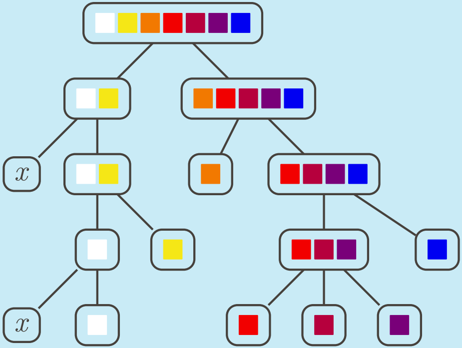
ii. $\mathfrak{P}$ is a non-identity echelonless path and $\mathfrak{Q}$ is one of the subpaths extracted from $\mathfrak{P}$.

We denote by $\leq$ to the reflexive transitive closure of $\prec$.

Proposition. $\leq$ is an Artinian order on $\mathrm{Pth}_{\mathcal{A}}$ whose minimal elements are identity paths and echelons.

DECOMPOSITION

Example



CATEGORIAL SIGNATURE

We define the **categorial signature** determined by the rewriting system $\boldsymbol{\mathcal{A}}$ to be the signature that enlarges Σ with

i. the rewriting rules in $\mathcal{A}$ as constants;
ii. two unary operations sc and tg;
iii. a binary operation $\circ$.

We will denote this signature with $\Sigma^{\boldsymbol{\mathcal{A}}}$.

THE CURRY-HOWARD MAPPING

The Curry-Howard mapping is defined by **Artinian recursion**

$$\mathrm{CH}\colon \mathrm{Pth}_{\boldsymbol{\mathcal{A}}} \longrightarrow \mathrm{T}_{\Sigma^{\boldsymbol{\mathcal{A}}}}(X)$$

1. For minimal paths

$$\mathrm{CH}(\mathrm{ip}(P)) = P; \qquad \mathrm{CH}(\mathrm{Ech}(\mathfrak{p})) = \mathfrak{p}.$$

2. For non-minimal paths

$$\mathrm{CH}(\mathfrak{P}) = \begin{cases} \mathrm{CH}(\mathfrak{P}^{i,|\mathfrak{P}|-1}) \circ \mathrm{CH}(\mathfrak{P}^{0,i-1}); \\ \sigma((\mathrm{CH}(\mathfrak{Q}_j))_{j\in n}). \end{cases}$$

THE CURRY-HOWARD MAPPING

Example

$$
\begin{aligned}
\mathfrak{P}:\ \oplus(x, \oplus(x, y)) &\to \oplus(x, \oplus(x, z)) \\
&\to \oplus(x, z) \\
&\to \odot(\boxdot(z, x), z, \boxdot(x, x)) \\
&\to \odot(\boxdot(z, x), x, \boxdot(x, x)) \\
&\to \odot(y, x, \boxdot(x, x)) \\
&\to \odot(y, x, z) \\
&\to \top
\end{aligned}
$$

$$
\mathrm{CH}(\mathfrak{P}) = ((\color{blue}\blacksquare\color{black} \circ (\odot(\color{purple}\blacksquare\color{black}, \color{red}\blacksquare\color{black}, \color{violet}\blacksquare\color{black}))) \circ \color{orange}\blacksquare\color{black}) \circ (\oplus(x, \color{yellow}\blacksquare\color{black} \circ \oplus(x, \color{white}\blacksquare\color{black})))
$$

THE ALGEBRA OF PATHS

Proposition. The set $\mathrm{Pth}_{\mathcal{A}}$ has structure of partial $\Sigma^{\mathcal{A}}$-algebra, that we will denote by $\mathbf{Pth}_{\mathcal{A}}$, where the operations are given by

$$\mathbf{sc}(\mathfrak{P}) = \mathrm{ip}(\mathrm{sc}(\mathfrak{P})); \qquad \mathbf{tg}(\mathfrak{P}) = \mathrm{ip}(\mathrm{tg}(\mathfrak{P}));$$

$$\boldsymbol{\mathfrak{p}} = \mathrm{Ech}(\mathfrak{p}); \qquad \mathfrak{Q} \boldsymbol{\circ} \mathfrak{P} = \mathfrak{Q} \circ \mathfrak{P}.$$

THE ALGEBRA OF PATHS

If $\sigma \in \Sigma_n$ and $(\mathfrak{P}_j)_{j\in n} \in \mathrm{Pth}^n_{\boldsymbol{\mathcal{A}}}$, then

$$\boldsymbol{\sigma}((\mathfrak{P}_j)_{j\in n}) : \quad \begin{array}{lccccccc}
\sigma(& \mathrm{sc}(\mathfrak{P}_0) & , & \mathrm{sc}(\mathfrak{P}_1) & , & \cdots & , & \mathrm{sc}(\mathfrak{P}_{n-1})\,) \\
 & \downarrow \mathfrak{P}_0 & & \| & & & & \| \\
\sigma(& \mathrm{tg}(\mathfrak{P}_0) & , & \mathrm{sc}(\mathfrak{P}_1) & , & \cdots & , & \vdots \quad) \\
 & \| & & \downarrow \mathfrak{P}_1 & & & & \| \\
\sigma(& \vdots & , & \mathrm{tg}(\mathfrak{P}_1) & , & \ddots & , & \mathrm{sc}(\mathfrak{P}_{n-1})\,) \\
 & \| & & \| & & & & \downarrow \mathfrak{P}_{n-1} \\
\sigma(& \mathrm{tg}(\mathfrak{P}_0) & , & \mathrm{tg}(\mathfrak{P}_1) & , & \cdots & , & \mathrm{tg}(\mathfrak{P}_{n-1})\,)
\end{array}$$

Proposition. $\sigma((\mathfrak{P}_j)_{j\in n})$ is an echelonless path.

THE KERNEL OF THE CURRY-HOWARD MAPPING

Proposition. CH is a Σ-homomorphism but not necessarily a $\Sigma^{\boldsymbol{\mathcal{A}}}$-homomorphism.

$$\mathrm{CH}(\mathrm{ip}(P)) = \mathrm{CH}(\mathrm{ip}(P) \circ \mathrm{ip}(P)) \neq \mathrm{CH}(\mathrm{ip}(P)) \circ \mathrm{CH}(\mathrm{ip}(P)).$$

Proposition. $\mathrm{Ker}(\mathrm{CH})$ is a closed $\Sigma^{\boldsymbol{\mathcal{A}}}$-congruence.

The quotient $\mathrm{Pth}_{\boldsymbol{\mathcal{A}}}/\mathrm{Ker}(\mathrm{CH})$ will be denoted by $[\mathrm{Pth}_{\boldsymbol{\mathcal{A}}}]$ and the class of a path $\mathfrak{P}$ will be denoted by $[\mathfrak{P}]$.

THE QUOTIENT OF PATHS

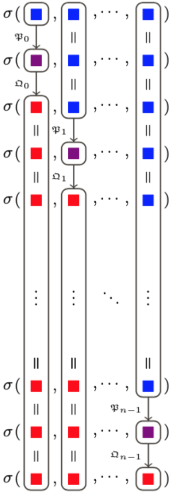
The quotient $[\mathrm{Pth}_{\boldsymbol{\mathcal{A}}}]$ has structure of **partial $\Sigma^{\boldsymbol{\mathcal{A}}}$-algebra, partially ordered set, and category**.

Furthermore, the operations $\sigma \in \Sigma$ of arity n are **functors** from $[\mathrm{Pth}_{\boldsymbol{\mathcal{A}}}]^n$ to $[\mathrm{Pth}_{\boldsymbol{\mathcal{A}}}]$, since

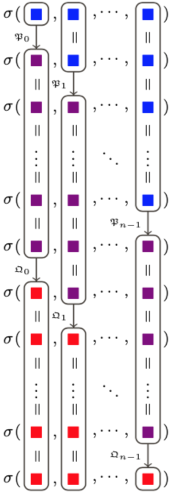
$$\mathbf{sc}\left(\sigma\left(([\mathfrak{P}_j])_{j\in n}\right)\right) = \sigma\left((\mathbf{sc}([\mathfrak{P}_j]))_{j\in n}\right)$$
$$\mathbf{tg}\left(\sigma\left(([\mathfrak{P}_j])_{j\in n}\right)\right) = \sigma\left((\mathbf{tg}([\mathfrak{P}_j]))_{j\in n}\right)$$
$$\sigma\left(([\mathfrak{Q}_j]\circ[\mathfrak{P}_j])_{j\in n}\right) = \sigma\left(([\mathfrak{Q}_j])_{j\in n}\right)\circ\sigma\left(([\mathfrak{P}_j])_{j\in n}\right)$$

This is a **categorial Σ-algebra** that we denote it by $[\mathbf{Pth}_{\boldsymbol{\mathcal{A}}}]$.

THE QUOTIENT OF PATHS



$\sigma((\mathfrak{Q}_j \circ \mathfrak{P}_j)_{j\in n})$



$\sigma((\mathfrak{Q}_j)_{j\in n}) \circ \sigma((\mathfrak{P}_j)_{j\in n})$

A CURRY-HOWARD RESULT

Theorem. There exists a pair of inverse mappings

$$[\mathbf{Pth}_{\mathcal{A}}] \underset{\mathrm{ip}^{\mathrm{fc}}}{\overset{\mathrm{CH}}{\rightleftarrows}} [\mathbf{PT}_{\mathcal{A}}] \quad (\cong)$$

- isomorphisms of partial $\Sigma^{\mathcal{A}}$-algebras;
- order isomorphisms;
- isomorphisms of categories.

SECOND-ORDER REWRITING SYSTEMS

This process can be **iterated**.

1. We introduce the notion of first-order translation T.
2. For every term class $[M] \in [\mathrm{PT}_{\mathcal{A}}]$, and every $M' \in [M]$.

$$[T(M)] = [T(M')].$$

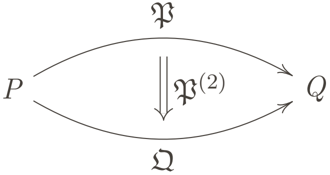
3. We introduce the notion of second-order rewriting rules as pairs $\mathfrak{p}^{(2)} = ([M], [N])$ with the condition

$$\mathrm{sc}\left(\mathrm{ip}^{\mathrm{fc}}(M)\right) = \mathrm{sc}\left(\mathrm{ip}^{\mathrm{fc}}(N)\right); \quad \mathrm{tg}\left(\mathrm{ip}^{\mathrm{fc}}(M)\right) = \mathrm{tg}\left(\mathrm{ip}^{\mathrm{fc}}(N)\right).$$

4. We introduce the notion of **second-order paths**.

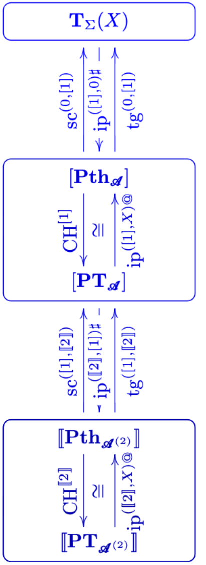
SECOND-ORDER PATHS

A second-order path $\mathfrak{P}^{(2)}$ has the form

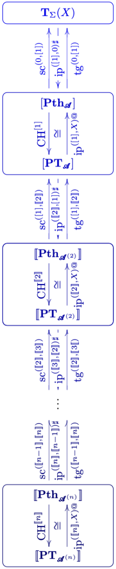


Mutatis mutandis **we recover the previous results**.

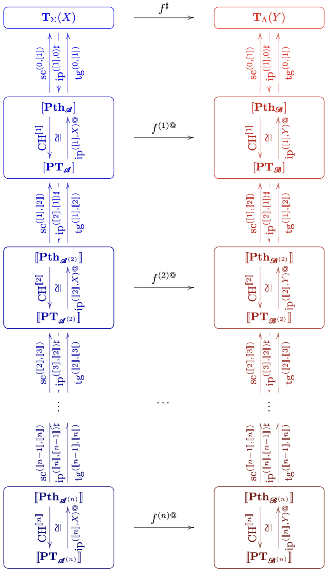
SECOND-ORDER RESULTS



N-TH ORDER RESULTS



MORPHISMS



$\mathbf{T}_\Sigma(X)$
$f^\sharp$
$\mathbf{T}_\Lambda(Y)$
$[\mathbf{Pth}_{\mathscr{A}}]$
$f^{(1)@}$
$[\mathbf{Pth}_{\mathscr{B}}]$
$[\mathbf{PT}_{\mathscr{A}}]$
$[\mathbf{PT}_{\mathscr{B}}]$
$[\![\mathbf{Pth}_{\mathscr{A}^{(2)}}]\!]$
$f^{(2)@}$
$[\![\mathbf{Pth}_{\mathscr{B}^{(2)}}]\!]$
$[\![\mathbf{PT}_{\mathscr{A}^{(2)}}]\!]$
$[\![\mathbf{PT}_{\mathscr{B}^{(2)}}]\!]$
$[\![\mathbf{Pth}_{\mathscr{A}^{(n)}}]\!]$
$f^{(n)@}$
$[\![\mathbf{Pth}_{\mathscr{B}^{(n)}}]\!]$
$[\![\mathbf{PT}_{\mathscr{A}^{(n)}}]\!]$
$[\![\mathbf{PT}_{\mathscr{B}^{(n)}}]\!]$

MORPHISMS

To determine a **morphism** from $\boldsymbol{\mathcal{A}}^{(n)}$ to $\boldsymbol{\mathcal{B}}^{(n)}$ we will assign

- to every variable in X a term in $\mathrm{T}_\Gamma(Y)$
- to every operation in Σ a **derived operation** in $\mathrm{T}_\Gamma(Y)$
- to every k-th rewriting rule in $\mathcal{A}^{(k)}$ a k-th order path in $\mathrm{Pth}_{\boldsymbol{\mathcal{B}}^{(k)}}$ respecting sources and targets

The final mapping $f^{(k)@}\colon [\![\mathbf{Pth}_{\boldsymbol{\mathcal{A}}^{(k)}}]\!] \longrightarrow [\![\mathbf{Pth}_{\boldsymbol{\mathcal{B}}^{(k)}}]\!]$, is obtained by **Artinian recursion** and by **universal property** on the quotients.

FUTURE WORK

1. Towers of rewriting systems.
2. Projective limits of rewriting systems.
3. Classifying spaces.

La possibilité de la traduction implique l’existence d’un invariant. Traduire, c’est précisément dégager cet invariant.

—H. Poincaré.

Thanks!